TWO DISCRETE INEQUALITIES OF GRÜSS TYPE VIA PÓLYA-SZEGÖ AND SHISHA RESULTS FOR REAL NUMBERS

S.S. DRAGOMIR AND L. KHAN

Abstract. Some new Grüss type discrete inequalities for nonnegative real numbers and applications for the moments of guessing mappings are given.

1. Introduction

In 1950, Biernacki, Pidek and Ryll-Nardzewski [1] proved the following Grüss type discrete inequality.

If \( \vec{a} = (a_1, \ldots, a_n) \) and \( \vec{b} = (b_1, \ldots, b_n) \) are such that there exists the real numbers \( a, A, b, B \) with

\[
\begin{align*}
(1.1) \quad a &\leq a_i \leq A, \quad b \leq b_i \leq B, \quad i \in \{1, \ldots, n\}
\end{align*}
\]

then

\[
(1.2) \quad \left| C_n (\vec{a}, \vec{b}) \right| \leq \frac{1}{n} \left[ \frac{n}{2} \right] \left( 1 - \frac{1}{n} \left[ \frac{n}{2} \right] \right) (A - a) (B - b)
\]

\[
= \frac{1}{n^2} \left[ \frac{n^2}{4} \right] (A - a) (B - b)
\]

\[
\leq \frac{1}{4} (A - a) (B - b)
\]

where

\[
C_n (\vec{a}, \vec{b}) := \frac{1}{n} \sum_{i=1}^{n} a_i b_i - \frac{1}{n} \sum_{i=1}^{n} a_i \cdot \frac{1}{n} \sum_{i=1}^{n} b_i.
\]

A weighted version of the above result has been obtained in 1988 by Andrica and Badea [2].

Let \( \vec{a}, \vec{b} \) satisfy (1.1) and \( \vec{p} = (p_1, \ldots, p_n) \) be an \( n \)-tuple of nonnegative numbers with \( P_n > 0 \). If \( S \) is a subset of \( \{1, \ldots, n\} \) that minimises the expression

\[
(1.3) \quad \left| \sum_{i \in S} p_i - \frac{1}{2} P_n \right|,
\]

then

\[
(1.4) \quad C_n (\vec{p}, \vec{a}, \vec{b}) \leq \frac{P_S}{P_n} \left( 1 - \frac{P_S}{P_n} \right) (A - a) (B - b)
\]

\[
\leq \frac{1}{4} (A - a) (B - b),
\]

Date: May 24, 2002.

1991 Mathematics Subject Classification. Primary 26D15; Secondary 94A05.

Key words and phrases. Integral Inequalities, Grüss Type Discrete Inequalities, Pólya-Szegö Inequality, Shisha Inequality.
where $P_S := \sum_{i \in S} p_i$ where

$$C_n(\bar{p}, \bar{a}, \bar{b}) := \frac{1}{P_n} \sum_{i=1}^{n} p_i a_i b_i - \frac{1}{P_n} \sum_{i=1}^{n} p_i a_i - \frac{1}{P_n} \sum_{i=1}^{n} p_i b_i.$$  

Recently, Dragomir and Booth \[3\] obtained the following result.

If $\bar{a}$, $\bar{b}$ are real $n$–tuples and $\bar{p}$ is nonnegative with $P_n > 0$, then

$$|C_n(\bar{p}, \bar{a}, \bar{b})| \leq \max_{1 \leq j \leq n-1} |\Delta a_j| \max_{1 \leq j \leq n-1} |\Delta b_j| C_n(\bar{p}, \bar{e}, \bar{e})$$  

where $\bar{e} = (1, 2, \ldots, n)$ and $\Delta a_j := a_{j+1} - a_j$ is the forward difference, and $j = 1, \ldots, n - 1$. Note that

$$C_n(\bar{p}, \bar{e}, \bar{e}) = \frac{1}{P_n} \sum_{i=1}^{n} i^2 p_i - \left(\frac{1}{P_n} \sum_{i=1}^{n} i p_i\right)^2.$$  

In particular, we have

$$|C_n(\bar{a}, \bar{b})| \leq \frac{1}{12} \left(n^2 - 1\right) \max_{1 \leq j \leq n-1} |\Delta a_j| \max_{1 \leq j \leq n-1} |\Delta b_j|.$$  

The constant $\frac{1}{12}$ is best possible.

In 2002, Dragomir \[4\] extended the above result for the $p$–norm. Namely, he proved that

$$|C_n(\bar{p}, \bar{a}, \bar{b})| \leq \frac{1}{P_n^2} \sum_{1 \leq i < j \leq n} (i - j) \left(\sum_{k=1}^{n-1} |\Delta a_k|^p\right)^{\frac{1}{p}} \left(\sum_{k=1}^{n-1} |\Delta b_k|^q\right)^{\frac{1}{q}}$$  

where $p > 1$, $\frac{1}{p} + \frac{1}{q} = 1$.

In particular, we have

$$|C_n(\bar{a}, \bar{b})| \leq \frac{1}{6} \cdot \frac{n^2 - 1}{n} \left(\sum_{k=1}^{n-1} |\Delta a_k|^p\right)^{\frac{1}{p}} \left(\sum_{k=1}^{n-1} |\Delta b_k|^q\right)^{\frac{1}{q}}.$$  

The constant $\frac{1}{6}$ is best possible.

The case of one-norm \[5\], can be stated as follows:

$$|C_n(\bar{p}, \bar{a}, \bar{b})| \leq \frac{1}{2} \cdot \frac{1}{P_n^2} \sum_{i=1}^{n} p_i (P_n - p_i) \sum_{k=1}^{n-1} |\Delta a_k| \sum_{k=1}^{n-1} |\Delta b_k|.$$  

In particular, we have

$$|C_n(\bar{a}, \bar{b})| \leq \frac{1}{2} \left(1 - \frac{1}{n}\right) \sum_{k=1}^{n-1} |\Delta a_k| \sum_{k=1}^{n-1} |\Delta b_k|.$$  

The constant $\frac{1}{2}$ is sharp.

Another direction was considered by Cerone and Dragomir in \[8\].

If $\bar{a}$, $\bar{b}$ are real $n$–tuples and $\bar{p}$ is a positive $n$–tuple and there exists $m, M \in \mathbb{R}$ such that

$$m \leq a_i \leq M,$$  

where $\bar{e} = (1, 2, \ldots, n)$.
then one has the inequality
\[ |C_n(\vec{p}, \vec{a}, \vec{b})| \leq \frac{1}{2} (M - m) \frac{1}{P_n} \sum_{i=1}^{n} p_i \left| b_i - \frac{1}{P_n} \sum_{j=1}^{n} p_j b_j \right|. \]

The constant \(\frac{1}{2}\) is best possible. In particular, we have
\[ |C_n(\bar{a}, \bar{b})| \leq \frac{1}{2} (M - m) \cdot \frac{1}{n} \sum_{i=1}^{n} a_i \frac{1}{n} \sum_{j=1}^{n} b_j. \]

The constant \(\frac{1}{2}\) is best possible.

In this paper we obtain different Grüss type discrete inequalities for nonnegative real numbers by the use of some counterpart results for the Cauchy-Buniakowski-Schwarz inequality. Application for the moments of guessing mapping are also given.

2. Discrete Inequalities

The following Grüss type inequality holds.

**Theorem 1.** Let \(\vec{a} = (a_1, \ldots, a_n)\) and \(\vec{b} = (b_1, \ldots, b_n)\) be two sequences of positive real numbers with
\[
0 < a_i \leq A < \infty \quad \text{and} \quad 0 < b_i \leq B < \infty \quad \text{for each} \quad i \in \{1, \ldots, n\}. 
\]

Then one has the inequality
\[
|C_n(\bar{a}, \bar{b})| \leq \frac{1}{4} \cdot \frac{(A - a) (B - b)}{\sqrt{AaBb}} \frac{1}{n} \sum_{i=1}^{n} a_i \cdot \frac{1}{n} \sum_{i=1}^{n} b_i.
\]

The constant \(\frac{1}{4}\) is best possible in (2.2) in the sense that it cannot be replaced by a smaller constant.

**Proof.** We have, by the Cauchy-Buniakowski-Schwarz inequality for double sums, the inequality
\[
|C_n(\bar{a}, \bar{b})| \\
\begin{align*}
= & \left| \frac{1}{2n^2} \sum_{i,j=1}^{n} (a_i - a_j)(b_i - b_j) \right| \\
\leq & \left| \frac{1}{2n^2} \sum_{i,j=1}^{n} |(a_i - a_j)(b_i - b_j)| \right| \\
\leq & \frac{1}{2n^2} \left[ \sum_{i,j=1}^{n} (a_i - a_j)^2 \sum_{i,j=1}^{n} (b_i - b_j)^2 \right]^{\frac{1}{2}} \\
= & \frac{1}{2n^2} \left[ 4 \left( \frac{n}{n} \sum_{i=1}^{n} a_i^2 - \left( \frac{n}{n} \sum_{i=1}^{n} a_i \right)^2 \right) \left( \frac{n}{n} \sum_{i=1}^{n} b_i^2 - \left( \frac{n}{n} \sum_{i=1}^{n} b_i \right)^2 \right) \right]^{\frac{1}{2}} \\
= & \left[ \frac{1}{n} \sum_{i=1}^{n} a_i^2 - \left( \frac{1}{n} \sum_{i=1}^{n} a_i \right)^2 \right]^{\frac{1}{2}} \left[ \frac{1}{n} \sum_{i=1}^{n} b_i^2 - \left( \frac{1}{n} \sum_{i=1}^{n} b_i \right)^2 \right]^{\frac{1}{2}}.
\end{align*}
\]
Utilising the Pólya-Szego inequality \[19\]

\[
1 \leq \frac{\sum_{i=1}^{n} z_i^2 \sum_{i=1}^{n} u_i^2}{(\sum_{i=1}^{n} z_i u_i)^2} \leq \frac{1}{4} \left( \sqrt{\frac{M_1 M_2}{m_1 m_2}} + \sqrt{\frac{m_1 m_2}{M_1 M_2}} \right)^2,
\]

provided \(0 < m_1 \leq z_i \leq M_1 < \infty, 0 < m_2 \leq u_i \leq M_2 < \infty, i \in \{1, \ldots, n\}\), we may state that

\[
n \sum_{i=1}^{n} a_i^2 \leq \frac{1}{4} \left( \sqrt{\frac{A}{a}} + \sqrt{\frac{a}{A}} \right)^2 = \frac{1}{4} \cdot (A + a)^2
\]
giving

\[
n \sum_{i=1}^{n} a_i^2 - \left( \sum_{i=1}^{n} a_i \right)^2 \leq \frac{1}{4} \cdot \frac{(A + a)^2}{a A} - 1 = \frac{(A - a)^2}{4 a A},
\]

that is,

\[
n \sum_{i=1}^{n} a_i^2 - \left( \sum_{i=1}^{n} a_i \right)^2 \leq \frac{(A - a)^2}{4 a A} \left( \sum_{i=1}^{n} a_i \right)^2.
\]

In a similar fashion, we obtain

\[
n \sum_{i=1}^{n} b_i^2 - \left( \sum_{i=1}^{n} b_i \right)^2 \leq \frac{(B - b)^2}{4 b B} \left( \sum_{i=1}^{n} b_i \right)^2.
\]

Using (2.3), (2.5) and (2.6), we deduce the desired inequality (2.2).

Now, assume that the inequality in (2.2) holds with a constant \(c > 0\), i.e.,

\[
|C_n(\bar{a}, \bar{b})| \leq \frac{(A - a)(B - b)}{\sqrt{a A b B}} \frac{1}{n} \sum_{i=1}^{n} a_i \cdot \frac{1}{n} \sum_{i=1}^{n} b_i.
\]

If we choose \(n = 2, a_1 = b_1, a_2 = b_2, a_1 = a, a_2 = A\), then from (2.7) we get

\[
\frac{1}{2} (a^2 + A^2) - \frac{1}{4} (a + A)^2 \leq c \frac{(A - a)^2}{a A} \cdot \frac{(a + A)^2}{4}
\]
giving

\[
\frac{1}{4} (A - a)^2 \leq c \frac{(A - a)^2}{a A} \cdot \frac{(a + A)^2}{4}
\]
from where we get

\[
(2.8) \quad aA \leq c(a + A)^2 \quad \text{for any } 0 < a < A < \infty.
\]

Let \(a = 1 - \varepsilon, A = 1 + \varepsilon, \) with \(\varepsilon \in (0, 1)\). Then from (2.8) we get \(1 - \varepsilon^2 \leq 4c\) for any \(\varepsilon \in (0, 1)\), which shows that \(c \geq \frac{1}{4}\).

\textbf{Remark 1.} We will now compare the inequality (2.2) with the Grüss inequality

\[
|C_n(\bar{a}, \bar{b})| \leq \frac{1}{4} (A - a)(B - b),
\]

providing \(a \leq a_i \leq A\) and \(b \leq b_i \leq B, i \in \{1, \ldots, n\}\).

We consider, for \(a, b > 0\), the quantity

\[
U := \frac{1}{\sqrt{a A b B}} \frac{1}{n} \sum_{i=1}^{n} a_i \cdot \frac{1}{n} \sum_{i=1}^{n} b_i
\]
and we will assume that \( a = b, A = B, a_i = b_i, i \in \{1, \ldots, n\} \). Thus

\[
U = \frac{(\sum_{i=1}^{n} a_i)^2}{n^2aA}.
\]

Choose \( n = 3, a_1 = a_2 = 1, a_3 = x \). Thus \( Aa = x \) and we have

\[
U(x) = \frac{(x+2)^2}{9x}.
\]

We observe that

\[
U(x) - 1 = \frac{x^2 - 9x + 4}{9x} = \frac{(x-1)(x-4)}{9x},
\]

showing that if \( x \in (0, 1] \cup [4, \infty) \), \( U(x) \geq 1 \) while for \( x \in (1, 4) \), \( U(x) < 1 \).

In conclusion, the bound provided by (2.2) is sometimes better, and at other times, worse than the bound provided by the Gr"uss inequality.

The second result of Gr"uss type is embodied in the following theorem.

**Theorem 2.** Let \( \bar{a} = (a_1, \ldots, a_n) \) and \( \bar{b} = (b_1, \ldots, b_n) \) be two sequences of positive real numbers satisfying (2.2). Then one has the inequality

\[
|C_n(\bar{a}, \bar{b})| \leq \left( \sqrt{A} - \sqrt{a} \right) \left( \sqrt{B} - \sqrt{b} \right) \sqrt{\frac{\sum_{i=1}^{n} a_i}{n} \cdot \frac{\sum_{i=1}^{n} b_i}{n}}.
\]

The constant \( c = 1 \) is the best possible in the sense that it cannot be replaced by a smaller constant.

**Proof.** We shall use Shisha’s inequality

\[
\sum_{i=1}^{n} z_i^2 - \frac{\sum_{i=1}^{n} z_i}{\sum_{i=1}^{n} y_i} \sum_{i=1}^{n} y_i^2 \leq \left( \sqrt{\frac{M_1}{m_2}} + \sqrt{\frac{m_1}{M_2}} \right)^2,
\]

provided \( 0 < m_1 \leq z_i \leq M_1 < \infty \) and \( 0 < m_2 \leq y_i \leq M_2 < \infty \).

If in (2.11) we choose \( z_i = a_i, y_i = 1 \), then we get

\[
0 \leq \frac{1}{n} \sum_{i=1}^{n} a_i^2 - \left( \frac{1}{n} \sum_{i=1}^{n} a_i \right)^2 \leq \frac{\sum_{i=1}^{n} a_i}{n} \left( \sqrt{A} - \sqrt{a} \right)^2.
\]

Similarly

\[
0 \leq \frac{1}{n} \sum_{i=1}^{n} b_i^2 - \left( \frac{1}{n} \sum_{i=1}^{n} b_i \right)^2 \leq \frac{\sum_{i=1}^{n} b_i}{n} \left( \sqrt{B} - \sqrt{b} \right)^2.
\]

Now, making use of (2.3), (2.12) and (2.13), we obtain the desired inequality (2.10).

To prove the sharpness of the constant, assume that (2.10) holds with a constant \( c > 0 \), i.e.,

\[
\left| \frac{1}{n} \sum_{i=1}^{n} a_i b_i - \frac{1}{n} \sum_{i=1}^{n} a_i \cdot \frac{1}{n} \sum_{i=1}^{n} b_i \right| \leq c \left( \sqrt{A} - \sqrt{a} \right) \left( \sqrt{B} - \sqrt{b} \right) \sqrt{\frac{\sum_{i=1}^{n} a_i}{n} \cdot \frac{\sum_{i=1}^{n} b_i}{n}}.
\]

If we choose \( n = 2, a_1 = b_1, a_2 = b_2, a_1 = a, a_2 = A \), then from (2.14) we get

\[
\frac{1}{4} (A - a)^2 \leq c \left( \sqrt{A} - \sqrt{a} \right)^2 \cdot \frac{a + A}{2},
\]
that is,
\[ \frac{1}{4} \left(\sqrt{A} - \sqrt{a}\right)^2 \left(\sqrt{A} + \sqrt{a}\right)^2 \leq c \left(\sqrt{A} - \sqrt{a}\right)^2 \cdot \frac{a + A}{2}, \]
giving for any $0 < a < A < \infty$ that
\[ (2.15) \left(\sqrt{A} + \sqrt{a}\right)^2 \leq 2c (a + A). \]

If in (2.15) we choose $a = 1 - \varepsilon$, $A = 1 + \varepsilon$, $\varepsilon \in (0, 1)$, we get $\left(\sqrt{1 - \varepsilon} + \sqrt{1 + \varepsilon}\right)^2 \leq 4c$. Letting $\varepsilon \to 0^+$, we deduce $c \geq 1$, and the theorem is proved. \[ \square \]

**Remark 2.** We shall show that at some times, the Grüss inequality (2.8) is better, and at other times, the inequality (2.10) is better.

If we choose $a_i = b_i$, $i = 1, n$, $a = b$, $A = B$, we have to compare
\[ I_1 := \frac{1}{4} (A - a)^2 \]
with
\[ I_2 := \left(\sqrt{A} - \sqrt{a}\right)^2 \sum_{i=1}^{n} a_i. \]
If we assume that $a = 0$, $A = 1$, then
\[ I_1 = \frac{1}{4}, \quad I_2 = \frac{\sum_{i=1}^{n} a_i}{n} \quad (i = 1, n) \]
showing that for $0 \leq a_i \leq 1$ with $\frac{\sum_{i=1}^{n} a_i}{n} < \frac{1}{4}$, (2.10) is better than the Grüss inequality while for $\frac{\sum_{i=1}^{n} a_i}{n} > \frac{1}{2}$, the Grüss inequality provides a better bound.

**Remark 3.** We will show now the fact that the bounds provided by (2.2) and (2.5) cannot generally be compared.

Assume that $a_i = b_i$, $(i = 1, \ldots, n)$, $a = b$, $A = b$ and consider
\[ J_1 := \frac{1}{4} \frac{(A - a)^2}{aA} \left(\frac{1}{n} \sum_{i=1}^{n} a_i\right)^2 \]
\[ J_2 := \left(\sqrt{A} - \sqrt{a}\right)^2 \frac{\sum_{i=1}^{n} a_i}{n}. \]
If we choose $a = 1$, $A = 4$, we get
\[ J_1 = \frac{9}{16} x^2, \quad J_2 = x \text{ where } x := \frac{\sum_{i=1}^{n} a_i}{n} \in [1, 4]. \]
We observe that $J_1 - J_2 = \frac{x(9x - 16)}{16}$ showing that for $x \in \left[1, \frac{16}{9}\right]$ the bound provided by (2.10) is better than the bound provided by (2.10) while for $x \in \left(\frac{16}{9}, 4\right)$, the conclusion is the other way around.

### 3. Applications for Moments of Guessing Mappings

In 1994, J.L. Massey [14] considered the problem of guessing the value taken on by a discrete random variable $X$ in one trial of a random experiment by asking questions of the form
\[ (3.1) \text{“Did } X \text{ take on its } i^{th} \text{ possible value?”} \]
until the answer is
\[ (3.2) \text{“Yes!”}. \]
This problem arises for instance when a cryptologist must try out possible secret keys one at a time after minimizing the possibilities by some cryptoanalysis.

Consider a random variable $X$ with finite range $X = \{x_1, \ldots, x_n\}$ and distribution $P_X(x_k) = p_k$ for $k = 1, 2, \ldots, n$.

A one-to-one function $G: \chi \to \{1, \ldots, n\}$ is a guessing function for $X$. Thus

\begin{equation}
E(G^m) := \sum_{k=1}^{n} k^m p_k
\end{equation}

is the $m$th moment of this function, provided we renumber the $x_i$ such that $x_k$ is always the $k$th guess.

In [14], Massey observed that, $E(G)$, the average number of guesses, is minimized by a guessing strategy that guesses the possible values of $X$ in decreasing order of probability.

In the same paper [14], Massey proved that

\begin{equation}
E(G) \geq \frac{1}{4} 2^{H(X)} + 1 \quad \text{provided} \quad H(X) \geq 2 \text{ bits},
\end{equation}

for an optimal guessing strategy, where $H(X)$ is the Shannon entropy

\begin{equation}
H(X) = - \sum_{i=1}^{n} p_i \log_2(p_i).
\end{equation}

He also has shown that $E(G)$ may be arbitrarily large when $H(X)$ is an arbitrarily small positive number such that there is no interesting upper bound on $E(G)$ in terms of $H(X)$.

In 1996, Arikan [15] has proved that any guessing algorithm for $X$ obeys the lower bound

\begin{equation}
E(G^\rho) \geq \left[\sum_{k=1}^{n} p_k^{\frac{1}{1+\rho}}\right]^{1+\rho}, \quad \rho \geq 0
\end{equation}

while an optimal guessing algorithm for $X$ satisfies

\begin{equation}
E(G^\rho) \leq \left[\sum_{k=1}^{n} p_k^{\frac{1}{1+\rho}}\right]^{1+\rho}, \quad \rho \geq 0.
\end{equation}

In 1997, Boztaş [16] proved that for $m \geq 1$, integer

\begin{equation}
E(G^m) \leq \frac{1}{m+1} \left[\sum_{k=1}^{n} p_k^{\frac{1}{1+m}}\right]^{1+m} + \frac{1}{m+1} \left\{ \left(\frac{m+1}{2}\right) E(G^{m-1}) - \left(\frac{m+1}{3}\right) E(G^{m-2}) + \cdots + (-1)^{m+1} \right\}
\end{equation}

provided the guessing strategy satisfies:

\begin{equation}
p_k^{\frac{1}{1+m}} \leq \frac{1}{k} \left( p_1^{\frac{1}{1+m}} + \cdots + p_k^{\frac{1}{1+m}} \right), \quad k = 1, \ldots, n-1.
\end{equation}

In 1997, Dragomir and Boztaş [17] obtained for any guessing sequence:

\begin{equation}
\left| E(G) - \frac{n+1}{2} \right| \leq \frac{(n-1)(n+1)}{6} \max_{1 \leq i < j \leq n} |p_i - p_j|,
\end{equation}
\[(3.11) \quad |E(G) - \frac{n+1}{2}| \leq \sqrt{\frac{(n-1)(n+1)(\|p\|_2^2 - 1)}{12}},\]

where \(\|p\|_2^2 = \sum_{i=1}^{n} p_i^2\) and

\[(3.12) \quad |E(G) - \frac{n+1}{2}| \leq \left[\frac{n+1}{2}\right] \left(n - \left[\frac{n+1}{2}\right]\right) \max_{1 \leq k \leq n} \left|p_k - \frac{1}{n}\right|,\]

where \([x]\) is the integer part of \(x\).

For other results on \(E(G^p)\), \(p > 0\) see also [18]. We mention only, by making use of Grüss inequality, one has for \(p, q > 0\) that

\[(3.13) \quad |E(G^{p+q}) - E(G^p)E(G^q)| \leq \frac{1}{4} (n^q - 1) (n^p - 1).\]

The above result may be complemented in the following way (see for example [11]).

**Theorem 3.** With the above assumptions, we have the inequality

\[(3.14) \quad |E(G^{p+q}) - \frac{1+n^q}{2} E(G^p) - \frac{1+n^p}{2} E(G^q) + \frac{1+n^q}{2} \cdot \frac{1+n^p}{2}| \leq \frac{1}{4} (n^q - 1) (n^p - 1).\]

for any \(p, q > 0\).

Applications for different particular instances of \(p, q > 0\) may be provided, but we omit the details.

The following result also holds [9].

**Theorem 4.** Assume \(S_n(p), p > 0\) denotes the sum of \(p\)-power of the first \(n\) natural numbers, that is

\[S_n(p) := \sum_{k=1}^{n} i^p.\]

If

\[p_i \leq (\geq) \frac{1}{n} \quad \text{for} \quad i \leq \left[\frac{S_n(p)}{n}\right]^{1/p}\]

and

\[p_i \geq (\leq) \frac{1}{n} \quad \text{for} \quad i \geq \left[\frac{S_n(p)}{n}\right]^{1/p} + 1\]

where \([x]\) denotes the integer part of \(x\), then we have the inequality

\[E(G^p) \geq (\leq) \frac{1}{n} S_n(p).\]

We are able now to state the first result for the moments of guessing mapping that may be obtained by the use of the inequality (2.2).

**Theorem 5.** If the probability distribution \((p_1, ..., p_n)\) satisfies the assumption

\[(3.15) \quad 0 < p_m \leq p_i \leq p_M \quad \text{for any} \quad i \in \{1, ..., n\},\]

then one has the inequality

\[|E(G^p) - \frac{1}{n} S_n(p)| \leq \frac{1}{4} \frac{(p_M - p_m)}{n^{p/2+1}} \cdot \frac{n^p - 1}{\sqrt{p_m p_M}} \cdot S_n(p).\]
In particular, for \( p = 1 \), we have the inequality
\[
\left| E(G) - \frac{n+1}{2} \right| \leq \frac{1}{8} \frac{(p_M - p_m)}{\sqrt{n}} \cdot \frac{n^2 - 1}{\sqrt{p_mpM}}.
\]

If one uses the other Grüss type inequality (2.10), then one may state the following result as well.

**Theorem 6.** If the probability distribution \((p_1, \ldots, p_n)\) satisfies the assumption (3.15), then one has the inequality
\[
\left| E(G^p) - \frac{1}{n} S_n(p) \right| \leq (\sqrt{p_M} - \sqrt{p_m}) (\sqrt{n^p} - 1) \sqrt{S_n(p)}.
\]

In particular, for \( p = 1 \), we have the inequality
\[
\left| E(G) - \frac{n+1}{2} \right| \leq (\sqrt{p_M} - \sqrt{p_m}) (\sqrt{n} - 1) \sqrt{\frac{n(n+1)}{2}}.
\]

**References**


School of Communications and Informatics, Victoria University of Technology, PO Box 14428, MCMC 8001, Victoria, Australia.

*E-mail address*: sever@matilda.vu.edu.au